Contents

- 1 Curie's advanced usage manual
- <u>2 Optimization</u>
 - 2.1 Compilation options
 - <u>2.1.1 Intel</u>
 - 2.1.1.1 Intel Sandy Bridge processors
 - <u>2.1.2 GNU</u>
- <u>3 Submission</u>

• 3.1 Choosing or excluding nodes

- <u>4 MPI</u>
 - 4.1 Embarrassingly parallel jobs and MPMD jobs
 - 4.2 BullxMPI
 - 4.2.1 MPMD jobs
 - 4.2.2 Tuning BullxMPI
 - 4.2.3 Optimizing with BullxMPI
 - 4.2.4 Debugging with BullxMPI
- 5 Process distribution, affinity and binding
 - 5.1 Introduction
 - 5.1.1 Hardware topology
 - 5.1.2 Definitions
 - <u>5.1.3 Process distribution</u>
 - 5.1.4 Why is affinity important for improving performance ?
 - 5.1.5 CPU affinity mask
 - <u>5.2 SLURM</u>
 - 5.2.1 Process distribution
 - 5.2.1.1 Curie hybrid node
 - 5.2.2 Process binding
 - 5.3 BullxMPI
 - <u>5.3.1 Process distribution</u>
 - 5.3.2 Process binding
 - <u>5.3.3 Manual process management</u>
- 6 Using GPU
 - 6.1 Two sequential GPU runs on a single hybrid node
- <u>7 Profiling</u>
 - <u>7.1 PAPI</u>
 - 7.2 VampirTrace/Vampir
 - <u>7.2.1 Basics</u>
 - 7.2.2 Tips
 - <u>7.2.3 Vampirserver</u>
 <u>7.2.4 CUDA profiling</u>
 - 7.3 Scalasca
 - 7.3.1 Standard utilization
 - 7.3.2 Scalasca + Vampir
 - 7.3.3 Scalasca + PAPI
 - 7.4 Paraver
 - 7.4.1 Trace generation
 - <u>7.4.2 Converting traces to Paraver format</u>
 - 7.4.3 Launching Paraver

Curie's advanced usage manual

If you have suggestions or remarks, please contact us : hotline.tgcc@cea.fr

Optimization

Compilation options

Compilers provides many options to optimize a code. These options are described in the following section.

Intel

- -opt_report : generates a report which describes the optimisation in stderr (-O3 required)
- -ip, -ipo : inter-procedural optimizations (mono and multi files). The command xiar must be used instead of ar to generate a static library file with objects compiled with -ipo option.
- -fast : default high optimisation level (-O3 -ipo -static). + Carefull : This option is not allowed using MPI, the MPI context needs to call some libraries which only exists in dynamic mode. This is incompatible with the -static option. You need to replace -fast by -O3 -ipo
- -ftz : considers all the denormalized numbers (like INF or NAN) as zeros at runtime.
- -fp-relaxed : mathematical optimisation functions. Leads to a small loss of accuracy.
- -pad : makes the modification of the memory positions operational (ifort only)

There are some options which allow to use specific instructions of Intel processors in order to optimize the code. These options are compatible with most of Intel processors. The compiler will try to generate these instructions if the processor allow it.

- -xSSE4.2 : May generate Intel® SSE4 Efficient Accelerated String and Text Processing instructions. May
 generate Intel® SSE4 Vectorizing Compiler and Media Accelerator, Intel® SSSE3, SSE3, SSE2, and SSE
 instructions.
- -xSSE4.1 : May generate Intel® SSE4 Vectorizing Compiler and Media Accelerator instructions for Intel processors. May generate Intel® SSSE3, SSE3, SSE2, and SSE instructions.
- -xSSSE3 : May generate Intel® SSSE3, SSE3, SSE2, and SSE instructions for Intel processors.
- -xSSE3 : May generate Intel® SSE3, SSE2, and SSE instructions for Intel processors.
- -xSSE2 : May generate Intel® SSE2 and SSE instructions for Intel processors.
- -xHost : this option will apply one of the previous options depending on the processor where the compilation is performed. This option is recommended for optimizing your code.

None of these options are used by default. The SSE instructions use the vectorization capability of Intel processors.

Intel Sandy Bridge processors

Curie thin nodes use the last Intel processors based on Sandy Bridge architecture. This architecture provides new vectorization instructions called AVX for Advanced Vector eXtensions. The option **-xAVX** allows to generate a specific code for Curie thin nodes.

Be careful, a code generated with -xAVX option runs only on Intel Sandy Bridge processors. Otherwise, you will get this error message:

Fatal Error: This program was not built to run in your system. Please verify that both the operating system and the processor support Inte I(R) AVX.

Curie login nodes are Curie large nodes with Nehalem-EX processors. AVX codes can be generated on these nodes through cross-compilation by adding -xAVX option. On Curie large node, the -xHost option will not generate a AVX code. If you need to compile with -xHost or if the installation requires some tests (like autotools/configure), you can submit a job which will compile on the Curie thin nodes.

GNU

There are some options which allow usage of specific set of instructions for Intel processors, in order to optimize code behavior. These options are compatible with most of Intel processors. The compiler will try to use these instructions if the processor allow it.

- -mmmx / -mno-mmx : Switch on or off the usage of said instruction set.
- -msse / -mno-sse : idem.
- -msse2 / -mno-sse2 : idem.
- -msse3 / -mno-sse3 : idem.
- -mssse3 / -mno-ssse3 : idem.
- -msse4.1 / -mno-sse4.1 : idem.
- -msse4.2 / -mno-sse4.2 : idem.
- -msse4 / -mno-sse4 : idem.
- -mavx / -mno-avx : idem, for Curie Thin nodes partition only.

Submission

Choosing or excluding nodes

SLURM provides the possibility to choose or exclude any nodes in the reservation for your job.

To choose nodes:

#!/bin/bash		
#MSUB-rMyjobPara		# Request name
#M5 UB -n 32	#	Number of tasks to use
#MS UB -T 1800	#	Ela psed time limit in se conds
#MSUB-oexample %I.o		# S ta nda rd output. % I is the job id
#MSUB-e example_%l.e		# Error output. % I is the job id
#MSUB-Apaxxxx		# Proje ct ID
#MS UB -E '-w curie [1000-1003]		# Include 4 nodes (curie 1000 to curie 1003)
set-x		
cd \${BRIDGE MS UB PWD}		
ccc mmm In out		

To exclude nodes:

#!/bin/bash		
#MSUB-rMyjobPara		# Request na me
#MS UB -n 32	#	Number of tasks to use
#MS UB -T 1800	#	Ela pse ditime limit in se conds
#MSUB-oexample %l.o		# S ta nda rd output. % I is the job id
#MSUB-e example_%l.e		# Error output. % I is the job id
#MSUB-Apaxxxx		# Proje ct ID
#MS UB -E '-x curie [1000-1003]'		# Exclude 4 nodes (curie 1000 to curie 1003)
set-x		
cd \${BRIDGE MS UB PWD}		
ccc mprun ./a .out		

MPI

Embarrassingly parallel jobs and MPMD jobs

- An embarrassingly parallel job is a job which launch independent processes. These processes need few or no communications
- A MPMD job is a parallel job which launch different executables over the processes. A MPMD job can be
 parallel with MPI and can do many communications.

These two concepts are separate but we present them together because the way to launch them on Curie is similar. An simple example in the Curie info page was already given.

In the following example, we use *ccc_mprun* to launch the job. *srun* can be used too. We want to launch *bin0* on the MPI rank 0, *bin1* on the MPI rank 1 and *bin2* on the MPI rank 2. We have first to write a shell script which describes the topology of our job:

launch_exe.sh:

We can then launch our job with 3 processes:

ccc_mprun -n 3 ./la unch_e xe .sh

The script *launch_exe.sh* must have execute permission. When *ccc_mprun* launches the job, it will initialize some environment variables. Among them, *SLURM_PROCID* defines the current MPI rank.

BullxMPI

MPMD jobs

BullxMPI (or OpenMPI) jobs can be launched with *mpirun* launcher. In this case, we have other ways to launch MPMD jobs (see embarrassingly parallel jobs section).

We take the same example in the embarrassingly parallel jobs section. There are then two ways for launching MPMD scripts

• We don't need the *launch_exe.sh* anymore. We can launch directly the job with *mpirun* command:

In the launch_exe.sh, we can replace SLURM_PROCID by OMPI_COMM_WORLD_RANK:

launch_exe.sh:

#!/bin/bash
if [\${OMPI_COMM_WORLD_RANK} -e q 0]
then
./bin0
TI
if [\${OMPLCOMM_WORLD_RANK} -e q 1]
the n
./bin1
fi
if [\${OMPLCOMM_WORLD_RANK} -e q 2]
the n
./bin2
fi

We can then launch our job with 3 processes:

mpirun -np 3 ./la un ch_e xe .sh

Tuning BullxMPI

BullxMPI is based on OpenMPI. It can be tuned with parameters. The command *ompi_info -a* gives you a list of all parameters and their descriptions.

ne suș ompi into -a	
MCA mpi: pa ra me te r "mpi_show_mca_pa ra ms" (curre nt va lue : < none >, da ta source : de fa ult va lue)	
Whe the r to show all MCA para me te r v a lue s during MP_INIT or not (good for re produca bility of MPI jobs for de bug purposes). Acce pte d v a lue s a re a II, de fa ult, file, a pi, a nd e nvironme nt	
- or a comma de limite d combina tion of the m	
.)	

Theses parameters can be modified with environment variables set before the *ccc_mprun* command. The form of the corresponding environment variable is *OMPI_MCA_xxxxx* where *xxxxx* is the parameter.

#1/hin/hash	
#MSIB_r Mulob Para	# Request name
#MS LIB -n 32	# Number of tasks to use
#MS LIB -T 1800	# Flansed time limit in seconds
#M511B-0 e xa mole %10	# Standard output % lis the job id
#MSUB-e example %/e	# Error output: % Lis the job id
#MS UB -A pa xxxx	# Project ID
set-x	
cd \${BRIDGE MS UB PWD}	
e xport OMPL MCA mpi show	mca params=a
ccc monin /a art	

Optimizing with BullxMPI

You can try theses parameters in order to optimize BullxMPI:

export OMPIMCA mpileave pinned=1

This setting improves the bandwidth for communication if the code uses the same buffers for communication during the execution.

export OMPI MCA btl openibuse eager rdma=1

This parameter optimizes the latence for short messages on Infiniband network. But the code will use more memory.

Be careful, theses parameters are not set by default. They can have influences on the behaviour of your codes.

Debugging with BullxMPI

Sometimes, BullxMPI codes can hang in any collective communication for large jobs. If you find yourself in this case, you can try this parameter:

e xport OMPI MCA coll="^ghc,tune d"

This setting disables optimized collective communications: it can slow down your code if it uses many collective operations.

Process distribution, affinity and binding

Introduction

Hardware topology

Machine (128GB)							
NUMANode P#0 (32G	iB)						
Socket P#0							
L3 (24MB)							
L2 (256KB)	L2 (256KB)	L2 (256KB)	L2 (256KB)	L2 (256KB)	L2 (256KB)	L2 (256KB)	L2 (256KB)
L1 (32KB)	L1 (32KB)	L1 (32KB)	L1 (32KB)	L1 (32KB)	L1 (32KB)	L1 (32KB)	L1 (32KB)
Core P#0	Core P#8	Core P#2	Core P#10	Core P#1	Core P#9	Core P#3	Core P#11
PU P#0	PU P#4	PU P#8	PU P#12	PU P#16	PU P#20	PU P#24	PU P#28
NUMANode P#1 (32G	3B)						
Socket P#2							
L3 (24MB)							
L2 (256KB)	L2 (256KB)	L2 (256KB)	L2 (256KB)	L2 (256KB)	L2 (256KB)	L2 (256KB)	L2 (256KB)
L1 (32KB)	L1 (32KB)	L1 (32KB)	L1 (32KB)	L1 (32KB)	L1 (32KB)	L1 (32KB)	L1 (32KB)
Core P#0	Core P#8	Core P#2	Core P#10	Core P#1	Core P#9	Core P#3	Core P#11
PU P#1	PU P#5	PU P#9	PU P#13	PU P#17	PU P#21	PU P#25	PU P#29
NUMANode P#2 (32G	6B)						
NUMANode P#2 (32G	B)						
NUMANode P#2 (32G Socket P#1 L3 (24MB)	iB)						
NUMANode P#2 (32G Sockel P#1 L3 (24MB) L2 (256KB)	L2 (258KB)	L2 (256KB)	L2 (258KB)	L2 (256KB)	L2 (256KB)	L2 (256KB)	L2 (250KB)
NUMANode P#2 (32G Socket P#1 L3 (24MB) L2 (256KB) L1 (32KB)	ав) L2 (256КВ) L1 (32КВ)	L2 (256KB) L1 (32KB)	L2 (250KB) L1 (32KB)	L2 (258KB) L1 (32KB)	L2 (250KB) L1 (32KB)	L2 (256KB) L1 (32KB)	L2 (250KB) L1 (32KB)
NUMANode P#2 (3203 Sockel P#1 L3 (24MB) L2 (250KB) L1 (32KB) Core P#0	L2 (259KB) L1 (32KB) Cone P#8	L2 (256KB) L1 (32KB) Come P#2	L2 (250KB) L1 (32KB) Come P#10	L2 (250KB) L1 (32KB) Come P#1	L2 (256KB) L1 (32KB) Core P#9	L2 (256KB) L1 (32KB) Come P#3	(2 (256KB) (1 (22KB) (1 (22KB) (0 cor P#11)
NUMANode P#2 (32G Socket P#1 L3 (24MB) L2 (256KB) L1 (32KB) Core P#0 PU P#2	ав) L2 (256КВ) L1 (32КВ) Сот Р#8 РU Р#6	L2 (256KB) L1 (32KB) Com P#2 PU P#10	L2 (259KB) L1 (22KB) Com P#10 PU P#14	L2 (259KB) L1 (32KB) Com P#1 PU P#18	L2 (259KB) L1 (32KB) Com P#9 PU P#22	L2 (250KB) L1 (22KB) Com P#3 PU P#20	L1 (256KB) L1 (2256KB) L1 (22KB) Com Pél11 PU Pe30
NUMANode P#2 (323 Socket P#1 L3 (24MB) L2 (256KB) L1 (32KB) L1 (32KB) Core P#0 PU P#2	ав) L2 (256КВ) L1 (22КВ) Соле Р#8 РU Р#6	L2 (256KB) L1 (32KB) Com P#2 PU P#10	L2 (250KB) L1 (32KB) Cow P#10 PU P#14	L2 (256KB) L1 (32KB) Com P#1 PU P#18	L2 (256KB) L1 (32KB) Core P#9 PU P#22	L2 (256KB) L1 (25KB) Cow P#3 PU P#20	L2 (256KB) L1 (22KS) Com P#11 PU P#30
NUMANode P#2 (3203 Sockel P#1 L3 (24MB) L2 (250KB) L1 (32KB) Core P#0 PU P#2 NUMANode P#3 (3203	L2 (259KB) L1 (32KB) Come P#8 PU P#6	L2 (256KB) L1 (32KB) Com P#2 PU P#10	L2 (250KB) L1 (32KB) Come P#10 PU P#14	L2 (259KB) L1 (32KB) Com P#1 PU P#18	L2 (250KB) L1 (32KB) Core P#9 PU P#22	L2 (256KB) L1 (32KB) Come P#3 PU P#26	(2 (256KB) (1 (22KB) (1 (22KB) Com P#11 PU P#30
NUMANode P#2 (32G Sockel P#1 L3 (24MB) L2 (256KB) L1 (32KB) Com P#0 PU P#2 NUMANode P#3 (32G Sockel P#3	12 (256KB) L1 (32KB) Com P#8 PU P#6 18)	L2 (256KB) L1 (32KB) Com P #2 PU P#10	L2 (250KB) L1 (32KB) Come P#10 PU P#14	L2 (256KB) L1 (22KB) Com P#1 PU P#18	L2 (256KB) L1 (32KB) Core P#9 PU P#22	L2 (256KB) L1 (22KB) Com P#3 PU P#26	L2 (256KB) L1 (32KB) Cow P#11 PU P#30
NUMANode P#2 (32G Sockel P#1 L3 (24MB) L1 (32KB) L1 (32KB) Com P#0 PU P#2 NUMANode P#3 (32G Sockel P#3 L3 (24MB)	38) L2 (256KB) L1 (32KB) Com P#8 PU P#6 38)	L2 (256KB) L1 (32KB) Com P#2 PU P#10	L2 (256KB) L1 (22KB) Cow P#10 PU P#14	L2 (256KB) L1 (32KB) Com P#1 PU P#18	L2 (259KB) L1 (32KB) Core P49 PU P#22	L2 (256KB) L1 (32KB) Cow P#3 PU P#28	L1 (20KB) L1 (27KB) Com P#11 PU P#30
NUMANode P#2 (320) Socket P#1 L3 (24MB) L2 (250KB) L1 (32KB) L1 (32KB) PU P#2 NUMANode P#3 (320) Socket P#3 L3 (24MB) L2 (250KB)	E2 (256KB) L1 (32KB) Come P#8 PU P#6 B) L2 (256KB)	L2 (256KB) L1 (22KB) Coxe P#2 PU P#10 L2 (256KB) L2 (256KB)	L2 (250KB) L1 (32KB) Come P#10 PU P#14 L2 (250KB)	L2 (256KB) L1 (32KB) Com P#I PU P#18 L2 (256KB)	L2 (250KB) L1 (32KB) Come P#9 PU P#22 L2 (250KB)	L2 (256KB) L1 (32KB) Com P#3 PU P#26 L2 (256KB)	(2 (256KB)) (1 (22KB)) (0 co P#11 PU P#30
NUMANode P#2 (32G Sockel P#1 L3 (24MB) L2 (256KB) L1 (32KB) Com P#0 PU P#2 NUMANode P#3 (32G Sockel P#3 L3 (24MB) L2 (256KB) L1 (32KB) L1 (32KB)	B) L2 (250KB) L1 (32KB) Com P#8 PU P#6 B) L2 (250KB) L2 (250KB) L1 (32KB) L1 (32KB)	L2 (256KB) L1 (32KB) Com P#2 PU P#10 L2 (256KB) L1 (32KB)	L2 (259KB) L1 (22KB) Com P#10 PU P#14 L2 (259KB) L1 (32KB)	L2 (259KB) L1 (32KB) Com P#1 PU P#18 L2 (259KB) L2 (259KB) L1 (32KB)	L2 (259KB) L1 (32KB) DU P#9 PU P#22 L2 (259KB) L1 (32KB)	L2 (250KB) L1 (22KB) Come P#3 PU P#20 L2 (250KB) L1 (22KB) L1 (32KB)	L1 (2250KB) L1 (32KB) Cow P#11 PU P#30 L1 (32KB) L1 (32KB) L1 (32KB)
NUMANode P#2 (323 Socket P#1 L3 (24MB) L1 (32KB) L1 (32KB) Du P#2 NUMANode P#3 (323 Socket P#3 L3 (24MB) L2 (256KB) L1 (32KB) Core P#0	IL2 (256КВ) L1 (22КВ) Соле Р#8 PU P#6 IL1 (22КВ) L2 (256КВ) L1 (22КВ) L1 (22КВ) L1 (22КВ) Cone P#8	L2 (256KB) L1 (32KB) Come P#2 PU P#10 L2 (256KB) L1 (32KB) L1 (32KB) Come P#2	L1 (256KB) L1 (32KB) Core P#10 PU P#14 L1 (22KB) L1 (22KB) L1 (22KB) Core P#10	L2 (256KB) L1 (32KB) Com P#1 PU P#19 L2 (256KB) L1 (32KB) L1 (32KB) Com P#1	L2 (256KB) L1 (32KB) Come P#9 PU P#22 L2 (256KB) L1 (32KB) L1 (32KB) Come P#9	L2 (256KB) L1 (32KB) Cow P#3 PU P#26 L1 (22KB) L1 (32KB) L1 (32KB) Cow P#3	L1 (226KB) L1 (22KB) Com P#11 L2 (256KB) L1 (22KB) L1 (22KB) L1 (22KB) Com P#11

Hardware topology of a Curie fat node

The hardware topology is the organization of cores, processors, sockets and memory in a node. The previous image was created with *hwloc*. You can have access to *hwloc* on Curie with the command *module load hwloc*.

Definitions

We define here some vocabulary:

- **Binding** : a Linux process can be bound (or stuck) to one or many cores. It means a process and its threads can run only on a given selection of cores. For example, a process which is bound to a socket on a Curie fat node can run on any of the 8 cores of a processor.
- Affinity : it represents the policy of resources management (cores and memory) for processes.
- **Distribution** : the distribution of MPI processes describes how theses processes are spread accross the core, sockets or nodes.

On Curie, the default behaviour for distribution, affinity and binding are managed by SLURM, precisely the *ccc_mprun* command.

Process distribution

We present here some example of MPI processes distributions.

• **block** or **round**: this is the standard distribution. From SLURM manpage: The block distribution method will distribute tasks to a node such that consecutive tasks share a node. For example, consider an allocation of two nodes each with 8 cores. A block distribution request will distribute those tasks to the nodes with tasks 0 to 7 on the first node, task 8 to 15 on the second node.



Block distribution by core

• **cyclic** by socket: from SLURM manpage, the cyclic distribution method will distribute tasks to a socket such that consecutive tasks are distributed over consecutive socket (in a round-robin fashion). For example, consider an allocation of two nodes each with 2 sockets each with 4 cores. A cyclic distribution by socket request will distribute those tasks to the socket with tasks 0,2,4,6 on the first socket, task 1,3,5,7 on the second socket. In the following image, the distribution is cyclic by socket and block by node.



Cyclic distribution by socket

• **cyclic** by node: from SLURM manpage, the cyclic distribution method will distribute tasks to a node such that consecutive tasks are distributed over consecutive nodes (in a round-robin fashion). For example, consider an allocation of two nodes each with 2 sockets each with 4 cores. A cyclic distribution by node request will distribute those tasks to the nodes with tasks 0,2,4,6,8,10,12,14 on the first node, task 1,3,5,7,9,11,13,15 on the second node. In the following image, the distribution is cyclic by node and block by socket.



Block distribution by node

Why is affinity important for improving performance?

Curie nodes are NUMA (Non-Uniform Memory Access) nodes. It means that it will take longer to access some regions of memory than others. This is due to the fact that all memory regions are not physically on the same bus.



NUMA node : Curie hybrid node

In this picture, we can see that if a data is in the memory module 0, a process running on the second socket like the 4th process will take more time to access the data. We can introduce the notion of *local data* vs *remote data*. In our example, if we consider a process running on the socket 0, a data is *local* if it is on the memory module 0. The data is *remote* if it is on the memory module 1.

We can then deduce the reasons why tuning the process affinity is important:

- Data locality improve performance. If your code use shared memory (like pthreads or OpenMP), the best choice is to regroup your threads on the same socket. The shared datas should be local to the socket and moreover, the datas will potentially stay on the processor's cache.
- System processes can interrupt your process running on a core. If your process is not bound to a core or to a socket, it can be moved to another core or to another socket. In this case, all datas for this process have to be moved with the process too and it can take some time.
- MPI communications are faster between processes which are on the same socket. If you know that two
 processes have many communications, you can bind them to the same socket.
- On Curie hybrid nodes, the GPUs are connected to buses which are local to socket. Processes can take longer time to access a GPU which is not connected to its socket.



NUMA node : Curie hybrid node with GPU

For all theses reasons, it is better to know the NUMA configuration of Curie nodes (fat, hybrid and thin). In the following section, we will present some ways to tune your processes affinity for your jobs.

CPU affinity mask

The affinity of a process is defined by a mask. A mask is a binary value which length is defined by the number of cores available on a node. By example, Curie hybrid nodes have 8 cores: the binary mask value will have 8 figures. Each figures will have 0 or 1. The process will run only on the core which have 1 as value. A binary mask must be read from right to left.

For example, a process which runs on the cores 0,4,6 and 7 will have as affinity binary mask: 11010001

SLURM and BullxMPI use theses masks but converted in hexadecimal number.

• To convert a binary value to hexadecimal:

\$ e cho "iba s e =2;oba s e =16;11010001"| bc 21202

• To convert a hexadecimal value to binary:

The numbering of the cores is the PU number from the output of hwloc.

SLURM

SLURM is the default launcher for jobs on Curie. SLURM manages the processes even for sequential jobs. We recommend you to use *ccc_mprun*. By default, SLURM binds processes to a core. The distribution is block by node and by core.

The option -E '--cpu bind=verbose' for ccc mprun gives you a report about the binding of processes before the run:

§ ccc. mpnun. E'-cpu_bind=verbose'-q hybrid-n.8./a.out cpu_bind=MASK - curie 7054, task 3 3 [3534]; mask 0x8set cpu_bind=MASK - curie 7054, task 0 0 [3531]; mask 0x8set cpu_bind=MASK - curie 7054, task 1 1 [3532]; mask 0x4 set cpu_bind=MASK - curie 7054, task 2 2 [3533]; mask 0x4 set cpu_bind=MASK - curie 7054, task 4 2 [3535]; mask 0x10 set cpu_bind=MASK - curie 7054, task 5 5 [3536]; mask 0x10 set cpu_bind=MASK - curie 7054, task 7 7 [3538]; mask 0x20 set cpu_bind=MASK - curie 7054, task 6 4 [3537]; mask 0x20 set cpu_bind=MASK - curie 7054, task 6 4 [3537]; mask 0x40 set

In this example, we can see the process 5 has 20 as hexadecimal mask or 00100000 as binary mask: the 5th process will run only on the core 5.

Process distribution

To change the default distribution of processes, you can use the option -*E* '-*m*' for *ccc_mprun*. With SLURM, you have two levels for process distribution: node and socket.

• Node block distribution:

ccc_mprun -E '-m block' ./a .out

• Node cyclic distribution:

ccc mprun - E'-m cyclic'./a.out

By default, the distribution over the socket is block. In the following examples for socket distribution, the node distribution will be block.

• Socket block distribution:

ccc_mprun -E '-m block:block' ./a .out

• Socket cyclic distribution:

ccc_mprun -E '-m block:cy clic' ./a .out

Curie hybrid node

On Curie hybrid node, each GPU is connected to a socket (see previous picture). It will take longer for a process to access a GPU if this process is not on the same socket of the GPU. By default, the distribution is block by core. Then the MPI rank 0 is located on the first socket and the MPI rank 1 is on the first socket too. The majority of GPU codes will assign GPU 0 to MPI rank 0 and GPU 1 to MPI rank 1. In this case, the bandwidth between MPI rank 1 and GPU 1 is not optimal.

If your code does this, in order to obtain the best performance, you should :

- use the block:cyclic distribution
- if you intend to use only 2 MPI processes per node, you can reserve 4 cores per process with the directive #MSUB -c 4. The two processes will be placed on two different sockets.

Process binding

By default, processes are bound to the core. For multi-threaded jobs, processes creates threads: these threads will be bound to the assigned core. To allow these threads to use other cores, SLURM provides the option -c to assign many cores to a process.

#!/bin/bash	
#MSUB-rMyjobPara	# Request na me
#M5 UB -n 8	# Number of tasks to use
#MS UB -c 4	#Assign 4 core s per process
#MS UB -T 1800	# Ela psed time limit in se conds
#MSUB-oexample_%I.o	# Standard output. % I is the job id
#MSUB-Apaxxxx	# Proje ct ID

In this example, our hybrid OpenMP/MPI code runs on 8 MPI processes and each process will use 4 OpenMP threads. We give here an example for the output with the verbose option for binding:

\$ ccc_mprun ./a .out			
cpu bind=MASK - curie 1139, task 5	5	5 [18761]: mask 0x40404040 set	
cpu bind=MASK - curie 1139, task 0	D	0 [18756]: mask 0x1010101 set	
cpu bind=MASK - curie 1139, task 1	1	1 [18757]: mask 0x10101010 set	
cpu bind=MASK - curie 1139, task 6	5	6 [18762]: mask 0x8080808 set	
cpu bind=MASK - curie 1139, ta sk 4	4	4 [18760]: mask 0x4040404 set	
cpu bind=MASK - curie 1139, task 3	3	3 [18759]: mask 0x20202020 set	
cpu bind=MAS K - curie 1139, ta sk 2	2	2 [18758]: mask 0x2020202 set	
cnu bind=MASK - curie 1139 task 7	7	7 [18763] mask 0x80808080 set	

We can see here the MPI rank 0 process is launched over the cores 0,8,16 and 24 of the node. These cores are all located on the node's first socket.

Remark: With the -c option, SLURM will try to gather at best the cores to have best performances. In the previous example, all the cores of a MPI process will be located on the same socket.

Another example:

(\$ ccc. mprun -n 1 -c 32 -E'cpu_bind=verbose'./a.out cpu_bind=MASK - curie 1017, task: 0 0 (34710); mask 0xffffffff set	
cpu_bind=MASK - curie 1017, task 0 0 [34710]: mask 0xfffffffset	

We can see the process is not bound to a core and can run over all cores of a node.

BullxMPI

#!/bin/bash

BullxMPI has its own process management policy. To use it, you have first to disable SLURM's process management policy by adding the directive #MSUB -E '--cpu_bind=none' . You can then use BullxMPI launcher mpirun:

#MSUB-rMyjobPara	# Request na me
#M5 UB -n 32	# Number of tasks to use
#MS UB -x	#Require a exclusive node
#MS UB -T 1800	# Ela psed time limit in se conds
#MSUB-oexample %1.o	# Standard output. % I is the job id
#MSUB-Apaxxxx	# Proje ct ID
#MS UB -E 'cpu_bind=none '	# Dis a ble de fa ult S LURM binding
mpirun -np 32 ./a .out	

Note: In this example, BullxMPI process management policy can be effective only on the 32 cores allocated by SLURM.

The default BullxMPI process management policy is:

- the processes are not bound
- the processes can run on all cores
- the default distribution is block by core and by node

The option --report-bindings gives you a report about the binding of processes before the run:

#:/DII/Dd 511	
#MSUB-rMyjobPara	# Request na me
#MS UB -n 32	# Number of tasks to use
#MS UB -x	#Require a exclusive node
#MS UB -T 1800	# Ela psed time limit in se conds
#MSUB-oexample %1.o	# Standard output. % I is the job id
#MSUB-Apaxxxx	# Proje ct ID
#MS UB -E 'cpu bind=none '	# Dis a ble de fa ult S LURM binding
mpirunre port-bindingsbind-	to-socketcpus-per-proc 4 -np 8 ./a .out

And there is the output:

+ mpirunbind-to-socketcpus-per-proc 4 -np 8 /a .out	
[curie 1342:19946] [[40080,0],0] odls :de fa ult:fork binding child [[40080,1],3] to socke t 1 cpus 2222222	
[curie 1342:19946] [[40080,0],0] odls:de fa ult:fork binding child [[40080,1],4] to socke t 2 cpus 4444444	
[[curie 1342:19946] [[40080,0],0] odls :de fa ult:fork binding child [[40080,1],5] to socke t 2 cpus 44444444	
[[curie 1342:19946] [[40080,0],0] odls :de fa ult:fork binding child [[40080,1],6] to socke t 3 cpus 88888888	
[[curie 1342:19946] [[40080,0],0] odls :de fa ult:fork binding child [[40080,1],7] to socke t 3 cpus 88888888	
[curie 1342:19946] [[40080,0],0] odls:de fa ult:fork binding child [[40080,1],0] to socke t 0 cpus 1111111	
[curie 1342:19946] [[40080,0],0] odls:de fa ult:fork binding child [[40080,1],1] to socke t 0 cpus 1111111	
[cune 1342:19946] [[40080,0],0] odls:de fa ult:fork binding child [[40080,1],2] to socket 1 cpus 2222222	

In the following paragraphs, we present the different possibilities of process distribution and binding. These options can be mixed (if possible).

Remark: the following examples use a whole Curie fat node. We reserve 32 cores with #MSUB -n 32 and #MSUB -x to have all the cores and to do what we want with them. This is only examples for simple cases. In others case, there may be conflicts with SLURM.

Process distribution



Cyclic distribution by socket:

#!/bin/bash	
#MSUB-rMyjobPara	# Request na me
#MS UB -n 32	# Number of tasks to use
#MS UB -x	#Require a exclusive node
#MS UB -T 1800	# Ela pse d time limit in se conds
#MSUB-oexample %1.o	# Standard output. % I is the job in
#MSUB-Apaxxxx	# Proje ct ID
#MS UB -E 'cpu bind=none '	# Dis a ble de fa ult S LURM binding
	, , , , , , , , , , , , , , , , , , ,
mpirunbysocket-np 32./a	out

Cyclic distribution by node:

#!/bin/bash	
#MSUB-rMyjobPara	# Requestname
#MS UB -n 32	# Number of tasks to use
#MS UB -N 16	
#MS UB -x	#Require exclusive nodes
#MS UB -T 1800	# Ela pse d time limit in se conds
#MSUB-oexample %1.o	# Standard output. % I is the job id
#MS UB - A pa xxxx	# Proje ct ID
#MS UB -E ' cpu bind=none '	# Dis a ble de fa ult S LURM binding
mpirunby node -np 32 ./a .or	.t

Process binding

No binding:

#!/bin/bash	
#MSUB-rMyjobPara	# Requestname
#M5 UB -n 32	# Number of tasks to use
#MS UB -x	#Require a exclusive node
#MS UB -T 1800	# Ela pse d time limit in se conds
#MSUB-oexample %1.o	# S ta nda rd output. % I is the job id
#MSUB-Apaxxxx	# Proje ct ID
#MS UB - E ' cpu bind=none '	# Disable de fault SLURM binding
. –	-
mpirunbind-to-none -np 32 .,	/a.out

Core binding:

#!/bin/bash	
#MSUB-rMyjobPara	# Request na me
#MS UB -n 32	# Number of tasks to use
#MS UB -x	#Require a exclusive node
#MS UB -T 1800	# Ela psed time limit in se conds
#MSUB-oexample %1.o	# Standard output. % I is the job id
#MSUB-Apaxxxx	# Proje ct ID
#MS UB -E 'cpu bind=none '	# Dis a ble de fa ult S LURM binding
mpirunbind-to-core -np 32 ./	a .out

Socket binding (the process and his threads can run on all cores of a socket):

#!/bin/bash	
#MSUB-rMyjob_Para	# Request na me
#MS UB -n 32	# Number of tasks to use
#MS UB -x	#Require a exclusive node
#MS UB -T 1800	# Ela pse ditime limit in se conds
#MSUB-oexample %1.o	# Standard output. % I is the job id
#MSUB-Apaxxxx	# Proje ct ID
#MS UB -E ' cpu bind=none '	# Dis a ble de fa ult S LURM binding
	-
mnin inhind-to-s ocke t -nn 32	/a out

You can specify the number of cores to assign to a MPI process:

#!/bin/bash	
#MSUB-rMyjobPara	# Request na me
#MS UB -n 32	# Number of tasks to use
#MS UB -x	#Require a exclusive node
#MS UB -T 1800	# Ela psed time limit in se conds
#MSUB-oexample_%I.o	# Standard output. % I is the job id
#MSUB-Apaxxxx	# Proje ct ID
#MS UB -E ' cpu bind=none '	# Dis a ble de fa ult S LURM binding
mpirunbind-to-s ocke tcous	-ne r-proc 4 -np 8 /a .out

Here we assign 4 cores per MPI process.

Manual process management

BullxMPI gives the possibility to manually assign your processes through a hostfile and a rankfile. An example:

#!/bin/bash	
#MSUB-r Myjob Para	# Request na me
#MS UB -n 32	# Number of tasks to use
#MS UB -x	#Require a exclusive node
#MS UB -T 1800	# Ela pse d time limit in se conds
#MSUB-oexample %I.o	# Standard output. % I is the job id
#MSUB-Apaxxxx	# Proje ct ID
#MS UB -E 'cpu bind=none '	# Dis a ble de fa ult S LURM binding

In this example, there are many steps :

- You have to create a *hostfile* here hostfile.txt where you put the hostname of all nodes your run will use
- You have to create a *rankfile* here rankfile.txt where you assign to each MPI rank the core where it can run. In our example, the process of rank 0 will have as affinity the core 0,1,2 and 3, etc... Be careful, the numbering of the core is different than the hwloc output: on Curie fat node, the eight first core are on the first socket 0, etc...
- you can launch *mpirun* by specifying the hostfile and the rankfile.

Using GPU

Two sequential GPU runs on a single hybrid node

To launch two separate sequential GPU runs on a single hybrid node, you have to set the environment variable CUDA_VISIBLE_DEVICES which enables GPUs wanted. First, create a script to launch binaries:

\$ cat la unch_e	e xe.sh
#!/bin/bash set-x	
e xport CUDA_V	VIS BLE_DEVICES = \${S LURM_PROCID} # the first process will see only the first GPU and the second process will see only the second GPU.
if [\$S LURM_PRC the n ./bin_1 > job_ fi if [\$S LURM_PRC the n ./bin_2 > job_ fi	ROCD -e q 0] b_\${S LURM_PROCD}.out ROCD -e q 1] b_\${S LURM_PROCD}.out

/!\ To work correctly, the two binaries have to been sequential (not using MPI).

Then run your script, making sure to submit two MPI processes with 4 cores per process:



So your first process will be located on the first CPU socket and the second process will be on the second CPU socket (each socket is linked with a GPU).

\$ ccc msub multijobs gpu.sh

Profiling

PAPI

PAPI is an API which allows you to retrieve hardware counters from the CPU. Here an example in Fortran to get the number of floating point operations of a matrix DAXPY:

program ma in implicit more include 190pa pi.ht ! inte ger, pa ra me ter :: size = 1000 inte ger, pa ra me ter :: ntime s = 10 double pre cision, dime ms insize , size) :: A,B,C inte ger, cise ison, dime ms insize , size) :: A,B,C inte ger, cine ness insize , size) :: A,B,C inte ger, cine ness insize , size) :: A,B,C inte ger, cine ness insize , size) :: A,B,C inte ger, cine ness insize , size) :: A,B,C inte ger, cine ness insize , size) :: A,B,C inte ger, cine ness insize , size) :: A,B,C inte ger, cine ness insize , size , size) :: A,B,C inte ger, cine ness insize , size , s end do ! Set up counters num events = 1 cal PAPf is tant, counters (event, num_events, retval) ! Clear the counter values cal PAPf read_counters (values, num_events, retval) ! DAXPf don=1,ntimes doi=1,size A(i,i) = 2.0*B(i,i) + C(i,i) end do end do end do ! Stop the counters and put the results in the array values cal PAPf is top, counters (values, num_events, retval) ! Print results f(event(1), BC, PAPI TOT_INS) the n print *, TOT Instructions: ', values (1) end if end print *, PP hs tructions: ', values (1) end print *, PP hs tructions: ', values (1) end if end print *, PP hs tructions: ', values (1) end if

To compile, you have to load the PAPI module :

ba sh-4.00 \$ module ba d pa pi/4.1.3 ba sh-4.00 \$ ifort-Is{PAPI_NC_DIR} pa pi.f90 \${PAPI_LIBS } ba sh-4.00 \$./a.out Number of ha rdwa re counters supported: 7 FP Instructions : 10046163

To get the available hardware counters, you can type "papi_avail" commande.

This library can retrieve the MFLOPS of a certain region of your code:

program main implicit none include '190pa pi.h' inte ger, para meter:: size = 1000 inte ger, para meter:: size = 1000 inte ger, para meter:: size = 100 double precision, dimension(size, size) :: A,B,C inte ger r:: tetval I Variable PAPI inte ger r:: retval re a (kind-4):: fipins I het PAPI retval = PAPI VER_CURRENT cal PAPI [brary_init; retval) if (retval.NEPAPI VER_CURRENT cal PAPI [brary_init; retval] end if do = 1.5 size do = 1.5

and the output:

bash-4.00 \$ module load pap/4.1.3
ba s h-4.00 \$ ifort -I\${PAPI_INC_DIR} pa pi_flops .f90 \${PAPI_LIBS }
bash-4.00 \$./a.out
Re a time : 6.1250001E-02
Proc_time : 5.1447589E-02
Tota flpins : 100056592
MFLOPS: 1944.826

If you want more precisions, you can contact us or visit PAPI website.

VampirTrace/Vampir

VampirTrace is a library which let you profile your parallel code by taking traces during the execution of the program. We present here an introduction of Vampir/Vampirtrace.

Basics

First, you must compile your code with VampirTrace compilers. In order to use VampirTrace, you need to load the vampirtrace module:

bash-4.00 \$ module load va mpirtra ce bash-4.00 \$ vtcc -c prog.c bash-4.00 \$ vtcc -o prog.e xe prog.o

Available compilers are :

- vtcc : C compiler
- vtc++, vtCC et vtcxx : C++ compilers
- vtf77 et vtf90 : Fortran compilers

To compile a MPI code, you should type :

bash-4.00 \$vtcc -vt:cc mpicc -g -c prog.c bash-4.00 \$vtcc -vt:cc mpicc -g -o prog.e xe prog.o

For others languages you have :

- vtcc -vt:cc mpicc : MPI C compiler
- vtc++ -vt:cxx mpic++, vtCC -vt:cxx mpiCC et vtcxx -vt:cxx mpicxx : MPI C++ compilers
- *vtf77 -vt:f77 mpif77* et *vtf90 -vt:f90 mpif90* : MPI Fortran compilers

By default, VampirTrace wrappers use Intel compilers. To change for another compiler, you can use the same method for MPI:

ba s h-4.00 \$ vtcc -vt:cc gcc -O2 -c prog.c ba s h-4.00 \$ vtcc -vt:cc gcc -O2 -o prog.e xe prog.o

To profile an OpenMP or a hybrid OpenMP/MPI application, you should add the corresponding OpenMP option for the compiler:

ba s h-4.00 \$ v tcc -ope nmp -O2 -c prog.c ba s h-4.00 \$ v tcc -ope nmp -O2 -o prog.e x e prog.o

Then you can submit your job. Here is an example of submission script:



At the end of execution, the program generates many profiling files :

ba s h-4.00 \$ ls a .out a .out.0.de f.z a .out.1.e ve nts .z ... a .out.otf

To visualize those files, you must load the vampir module:



Vampir window

If you need more information, you can contact us.

Vampirtrace allocate a buffer to store its profiling information. If the buffer is full, Vampirtrace will flush the buffer on disk. By default, the size of this buffer is 32MB per process and the maximum number of flushes is only one time. You can increase (or reduce) the size of the buffer: your code will also use more memory. To change the size, you have to initialize an environment variable :

e x port VT_BUFFER_S IZ E=64M ccc mprun ./prog.e x e

In this example, the buffer is set to 64 MB. We can increase the maximum number of flushes:

e xport VT_MAX_FLUS HES =10 ccc_mprun ./prog.e x e

If the value for VT_MAX_FLUSHES is 0, the number of flushes is unlimited.

By default, Vampirtrace will first store profiling information in a local directory (/tmp) of process. These files can be very large and fill the directory. You have to change this local directory with another location:

e x port VT PFORM LDIR=\$S CRATCHDIR

There are more Vampirtrace variables which can be used. See User Manual for more precisions.

Vampirserver

Traces generated by Vampirtrace can be very large: Vampir can be very slow if you want to visualize these traces. Vampir provides Vampirserver: it is a parallel program which uses CPU computing to accelerate Vampir visualization. Firstly, you have to submit a job which will launch Vampirserver on Curie nodes:

scativa mpirserver.sh #/bit/bash #/bit/bash #MSUB-rvampirserver # Request name #MSUB-raz # Number of tasks to use #MSUB-ration # Eapsed time limit in seconds #MSUB-ovampirserver_%Lo #MSUB-evampirserver_%Le #module load vampir \$ ccc.mspunvingd \$ module load vampir \$ ccc.msubvampirserver.sh

When the job is running, you will obtain this ouput:

```
$ ccc_mpp
USER_ACCOUNT_BATCHID_NCPU_QUEUE_PRIORITY_STATE_RLIM_RUN/START_SUSP_OLD_NAME___NODES
toto___genXXX_234481_32_large_210332_RUN_30.0m_1.3m__-__1.3m__vampirserver_curie1352
$ ccc_mpeek 234481
Found itce rise file : /usr/loca /vampir-7.3/bin/lic.da t
Rumring 31 ana lysis processes... (a bort with Ctrl-C or vngd-shutdown)
S erver listens on: curie 1352:30000
```

In our example, the Vampirserver master node is on curie1352. The port to connect is 30000. Then you can launch Vampir on front node. Instead of clicking on *Open*, you will click on *Remote Open*:

¥ Vampir (sur curie50)	
File Window Help	_
Connect to Server (sur curie50)	
Description Server Port	
Better Conce	

Connecting to Vampirserver

Fill the server and the port. You will be connected to vampirserver. Then you can open an OTF files and visualize it.

Notes:

- You can ask any number of processors you want: it will be faster if your profiling files are big. But be careful, it consumes your computing times.
- Don't forget to delete the Vampirserver job after your analyze.

CUDA profiling

Vampirtrace can collect profiling data from CUDA programs. As previously, you have to replace compilers by Vampirtrace wrappers. NVCC compiler should be replaced by *vtnvcc*. Then, when you run your program, you have to set an environment variable:

```
e xport e xport VT_CUDARTTRACE= y e s
ccc_mprun ./prog.e x e
```

Scalasca

Scalasca is a set of software which let you profile your parallel code by taking traces during the execution of the program. This software is a kind of parallel gprof with more information. We present here an introduction of Scalasca.

Standard utilization

First, you must compile your code by adding Scalasca tool before your call of the compiler. In order to use Scalasca, you need to load the scalasca module:

```
bash-4.00 $module load scalasca
bash-4.00 $scalasca -instrument mpicc-c prog.c
bash-4.00 $scalasca -instrument mpicc-o prog.exe prog.o
```

or for Fortran :

You can compile for OpenMP programs:

ba s h-4.00 \$ s ca la s ca -instrume nti ifort -ope nmp -c prog.f90 ba s h-4.00 \$ s ca la s ca -instrume nti ifort -ope nmp -o prog.e x e prog.o

You can profile hybrid programs:

```
ba sh-4.00 $ sca la sca -instrume nt mpif90 -ope nmp -O3 -c prog.f90
ba sh-4.00 $ sca la sca -instrume nt mpif90 -ope nmp -O3 -o prog.e xe prog.o
```

Then you can submit your job. Here is an example of submission script:

# !/onn/bash #M5UB-r Myjob_Para #M5UB-n 32 #M5UB-T 1800 #M5UB-o e xample_%1.o #M5UB-e e xample_%1.e	#Request name #Number of tasks touse #Bapsed time limit in seconds #Standard output. %I is the job id #Enroroutput. %I is the job id	
set-x cd\${BRIDGE_MSUB_PWD}		
export SCAN_MPILAUNCHER=ccc_mprun scala sca -a na lyze ccc_mprun ./prog.exe		

At the end of execution, the program generates a directory which contains the profiling files :

bash-4.00 \$ls epik_*

To visualize those files, you can type:



Scalasca

If you need more information, you can contact us.

Scalasca + Vampir

Scalasca can generate OTF tracefile in order visualize it with Vampir. To activate traces, you can add -t option to scalasca when you launch the run. Here is the previous modified script:

#I/bin/bash #MSUB-rMy/ob_Para #MSUB-n32 #MSUB-T1800 #MSUB-oexample_%Lo #MSUB-eexample_%Lo

Request name # Number of tasks touse # Bapsed time limit in seconds # Standardoutput. % is the job id # Ernoroutput. % is the job id

set-x cd\${BRIDGE_MSUB_PWD} At the end of execution, the program generates a directory which contains the profiling files :

bash-4.00 \$ls epik_*

To visualize those files, you can visualize them as previously. To generate the OTF trace files, you can type:

bash-4.00 \$ls epik_* bash-4.00 \$elg2otfepik_*

It will generate an OTF file under the epik * directory. To visualize it, you can load Vampir:

bash-4.00 \$ module load v a mpir bash-4.00 \$ v a mpir e pik_*/a.otf

Scalasca + PAPI

Scalasca can retrieve the hardware counter with PAPI. For example, if you want retrieve the number of floating point operations :

#/bin/bash #K5UB-rNvlobPana #Requestname	l
#MS-UB-r32 # Number of tasks to use	l
MPS US-1 LOOU # Da Jose U UITHE III NE I I SE CU UILE III NE I SE CU UILE III NE I CU UILE II ILI NE I CU UILE III NE I CU UILE II NE I CU UILE III NE I CU UILE III NE I CU UILE III NE I CU UIL	l
#MSUB-e example_%Le #Error output. % is the jobid	l
se t -x d \${BRDGE_M5 UB_PWD}	
export EPK_METRICS = PAPI_FP_OPS scala s ca - a na lyze mpirun ./prog.e xe	
	Î

Then the number of floating point operations will appear on the profile when you visualize it. You can retrieve only 3 hardware counters at the same time on Curie. The the syntax is:

e xport EPK_METRICS = "PAPI_FP_OPS : PAPI_TOT_CYC"

Paraver

Paraver is a flexible performance visualization and analysis tool that can be used to analyze MPI, OpenMP, MPI+OpenMP, hardware counters profile, Operating system activity and many other things you may think of!

In order to use Paraver tools, you need to load the paraver module:



Trace generation

The simpliest way to activate mpi instrumentation of your code is to dynamically load the library before execution. This can be done by adding the following line to your submission script:

e x port LD PRELOAD=\$LD PRELOAD:\$MPI TRACE LIBS

The instrumentation process is managed by Extrae and also need a configuration file in xml format. You will have to add next line to your submission script.

e xport EXTRAE_CONFIG_FILE=./e xtra e _config_file .xml

All detailled about how to write a config file are available in Extrae's manual which you can reach at \$EXTRAE_HOME/doc/user-guide.pdf. You will also find many examples of scripts in \$EXTRAE_HOME/examples/LINUX file tree.

You can also add some manual instrumentation in your code to add some specific user event. This is mandatory if you want to see your own functions in Paraver timelines.

If trace generation succeed during computation, you'll find a directory *set-0* containing some *.mpit* files in your working directory. You will also find a *TRACE.mpits* file which lists all these files.

Converting traces to Paraver format

Extrae provides a tool named mpi2prv to convert mpit files into a .prv which will be read by Paraver. Since it can be a long operation, we recommend you to use the parallel version of this tool, mpimpi2prv. You will need less processes than previously used to compute. An example script is provided below:

bash 4.00%, cat rebuildsh #M5UB-n marge #M5UB-n 8 #M5UB-T 1800 set-x cd \$#RIDEE_M5UB_PWD ccc_mprun mpimp2pvi-syn -e path_to_your_binany -f TRACE.mpts -ofile_to_be_analysed.prv

Launching Paraver

You just now have to launch "paraver file_to_be_analysed.prv". As Paraver may ask for high memory & CPU usage, it may be better to launch it through a submission script (do not forget then to activate the -X option in ccc_msub). For analyzing your data you will need some configurations files available in Paraver's browser under \$PARAVER_HOME/cfgs directory.



Paraver window